



渕一博記念コロキウム
『論理と推論技術：四半世紀の展開』

非単調性と帰納論理を取り入れたことで
論理プログラミングはどう変わったか？

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Incorporating Nonmonotonic and Inductive Inference into Logic Programming

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Thanks to:

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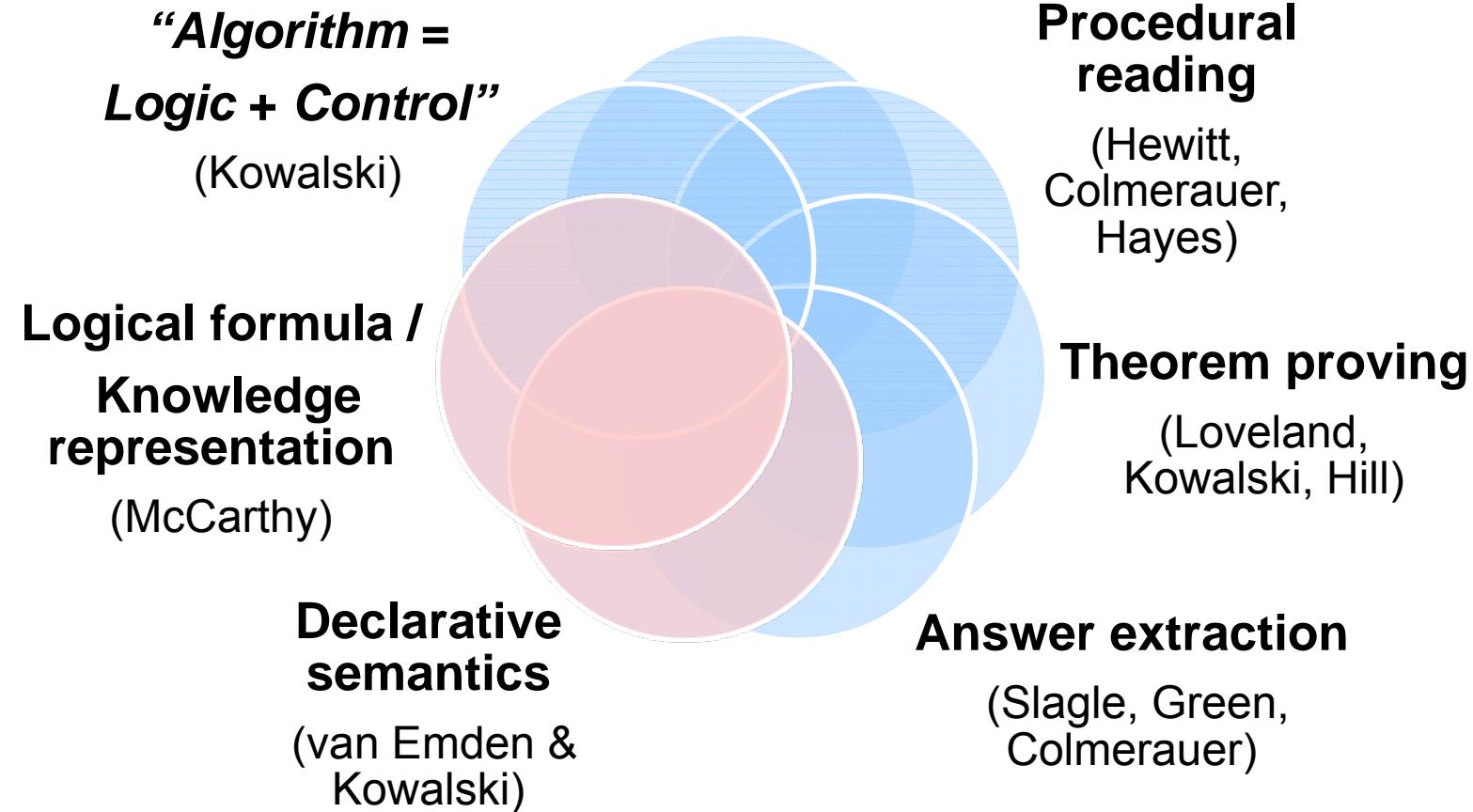
History of Thinking Machines

- 1936. Turing Machine
- 1943. McCulloch & Pitts: Boolean circuit model of brain
- 1950. Turing: “Computing Machinery and Intelligence”
- 1956. Dartmouth meeting: “Artificial Intelligence” adopted
- 1958. McCarthy: “Programs with Common Sense”
- 1950s. Early AI programs, including Samuel's checkers program, Newell & Simon's Logic Theorist
- 1960. McCarthy: LISP
- 1965. Robinson: complete algorithm for logical reasoning
- 1969. McCarthy & Hayes: situation calculus, frame problem
- 1971. Colmerauer & Kowalski: Prolog
- 1970s. Early development of knowledge-based systems
- 1980. McCarthy: circumscription
- 1982. Japan's Fifth Generation Computer Project**
- 1986. Neural networks return to popularity
- 1995. Emergence of intelligent agents
- 1997. Deep Blue

History of Logic Programming

- 1960. Davis & Putnum: testing satisfiability of propositional formulas
- 1962. Davis, Logemann & Loveland: improvement of DP → SAT
- 1965. Robinson: resolution principle
- 1968. Loveland: Model Elimination
- 1971. Kowalski & Kuehner: SL-resolution
- 1972. Colmerauer & Roussel: “*Programmation en Logique*”
- 1972. Kowalski: “The predicate calculus as a programming language”
- 1973. Hill: LUSH (SLD-)resolution
- 1974. van Emden & Kowalski: Scott’s fixpoint & Tarskian semantics
- 1977. Warren (David): Edinburgh Prolog compiler
- 1978. Clark: negation as failure (predicate completion)
- 1982. **Japan's Fifth Generation Computer Project**
- 1980s. Shapiro, Clark, Ueda: Concurrent Logic Programming
- 1987. Jaffer & Lassez: Constraint Logic Programming
- 1988. Gelfond & Lifschitz: stable model semantics
- 1991. Muggleton: “Inductive Logic Programming”
- 1992. Kakas, Kowalski & Toni: “Abductive Logic Programming”
- 1990s. Poole, Sato: Probabilistic Logic Programming

Interpreting Horn Logic Programs

$$H \ :- \ B_1, \dots, \ B_n.$$


Extending Horn Logic Programs

$H_1; \dots; H_m; \text{not } H_{m+1}; \dots; \text{not } H_n$

$\leftarrow B_1, \dots, B_m, \text{not } B_{m+1}, \dots, \text{not } B_n$

- Definite (Horn) program (H, B : atom)
- Normal logic program
- Extended logic program (H, B : literal)
- Extended disjunctive program
- General extended disjunctive program
- Nested program (H, B : rule)

Logic Programming and Nonmonotonic Reasoning

- 1969. McCarthy & Hayes: frame problem
- 1978. Reiter: closed-world assumption
- 1979. Clark: negation as failure (Compl)
- 1980. McCarthy: circumscription (Circ)
- 1980. Reiter: default logic (DL)
- 1982. Reiter: Circ implies Compl (sometimes)
- 1985. Moore: autoepistemic logic (AEL)
- 1987. Gelfond & Lifschitz: stratified LP as prioritized Circ
- 1988. Gelfond & Lifschitz: stable model semantics (LP as AEL)
- 1990. Gelfond & Lifschitz: answer set semantics (LP as DL)
- 1992. Inoue, Koshimura & Hasegawa: answer set computation in KL1
- 1996. Niemelä & Simons: smodels
- 1997. Eiter, Faber, Leone & Pfeifer: DLV
- 2007. Lifschitz, Lin: 2nd-order formalization of stable models (LP as Circ)

Answer Set Programming

A program is regarded as the **constraints** to be satisfied by **solutions**. Each solution is obtained by computing an **answer set (stable model)** of the program.

A program may have no, one, or multiple answer sets.

Program: $p ; \text{not } p \leftarrow,$
 $q ; \text{not } q \leftarrow,$
 $r ; \text{not } r \leftarrow,$
 $s \leftarrow p, r,$
 $s \leftarrow p, q, \text{not } r,$
 $\leftarrow \text{not } s.$

} generator
} tester
} goal

Answer Sets: $\{p, q, r, s\}, \{p, r, s\}, \{p, q, s\}.$

Answer Set Programming

A program is regarded as the **constraints** to be satisfied by **solutions**. Each solution is obtained by computing an **answer set (stable model)** of the program.

A program may have no, one, or multiple answer sets.

Program:

$$\begin{array}{l} p \leftarrow \mathbf{not} \ q, \mathbf{not} \ r, \\ q \leftarrow \mathbf{not} \ p, \mathbf{not} \ r, \\ r \leftarrow \mathbf{not} \ p, \mathbf{not} \ q, \\ s \leftarrow p, r, \\ s \leftarrow q, \mathbf{not} \ r, \\ \quad \leftarrow \mathbf{not} \ s. \end{array}$$

} generator
(exclusive
choice)
} tester
} goal

Answer Sets:

{q, s}.

Nonmonotonic Reasoning in ICOT

- 1984 Kitakami, Kunifugi, Miyachi & Furukawa: knowledge assimilation (SLP'84)
- 1986 Kunifugi, Tsurumaki & Furukawa: hypothesis-based reasoning
- 1986 Goebel, Poole & Furukawa: theory formation (ICLP'86)
- 1988 Ishizuka: *Hypothetical Reasoning WG*
- 1988 Arima, Satoh: work on circumscription (FGCS'88)
- 1989 Sakama: possible model semantics for disjunctive programs
- 1989 Nicolas Helft & David Poole in ICOT
- 1990 Inoue & Helft: circumscriptive theorem prover
- 1991 Helft, Inoue & Poole: QA in circumscription (IJCAI-91)
- 1991 Mark Stickel in ICOT
- 1991 Inoue: SOL-resolution (IJCAI-91)
- 1991 Inoue, Satoh & Iwayama: abduction in answer set semantics (ICLP'91)
- 1992 Inoue, Koshimura & Hasegawa: answer set comp. on MGTP (CADE'92)
- 1992 Ohta & Inoue: parallel abductive reasoning on MGTP & ATMS (FGCS'92)
- 1993 Inoue, Ohta, Hasegawa & Nakashima: abduction on MGTP (IJCAI-93)
- 1993 Inoue & Sakama: fixpoint semantics for abductive programs (ICLP'93)

Inoue & Sakama or Sakama & Inoue

1993. ICLP'93
1994.
1995.
1996. JLP
1997.
1998. KR'98
1999. AMAI
2000.
2001.
2002. ICLP'02
2003.
2004. JELIA'04
2005. CLIMA-05
2006.
2007.
2008. ACM.TOCL

ICLP'93

ICLP'93

JAR

ICLP'95

JICSLP'96

JLP

KR'98

AMAI

AIJ

LPNMR'93

KR'94

JLP

LPNMR'99

TPLP

IJCAI-95

ICLP'94



JLP

LPNMR'95 JLC

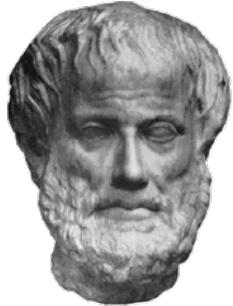
IJCAI-99

ILP'05
CLIMA-06
ICLP'06
AAAI-07

CLIMA-05
LJ.IGPL
AAMAS'07
LOPSTR'06

Abduction, Induction and Deduction

Analytic Reasoning



{ **Deduction**
(LP)

{ **consequence:**
from *prior knowledge*
to *necessary implications*

Synthetic Reasoning



{ **Induction**
(ILP)

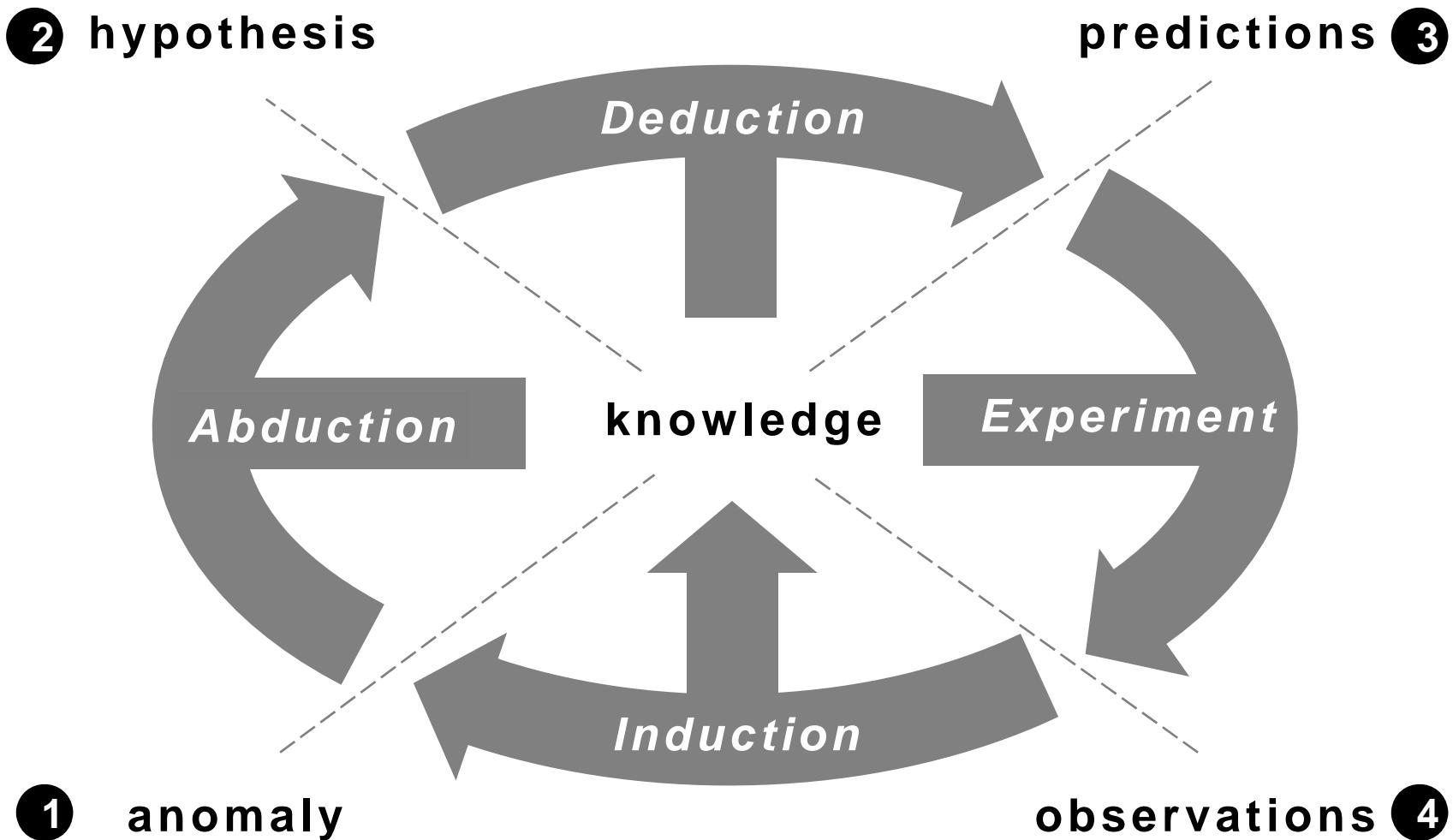
{ **generalization:**
from *observed samples*
to *wider populations*

{ **Abduction**
(ALP)

{ **explanation:**
from *given effects*
to *possible causes*

Scientific Discovery

Scientific Knowledge Development



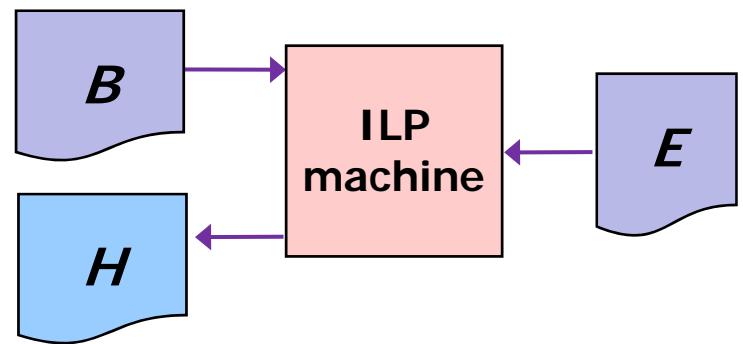
Abduction and Induction: Logical Framework

Input:

- B : background theory
- E : (positive) examples / observations

Output:

- H : hypothesis satisfying that
 - $B \wedge H \models E$
 - $B \wedge H$ is consistent.



Inverse Entailment (IE)

Computing a hypothesis H can be done **deductively** by:

$$B \wedge \neg E \models \neg H$$

We use a **consequence finding** technique for IE computation.

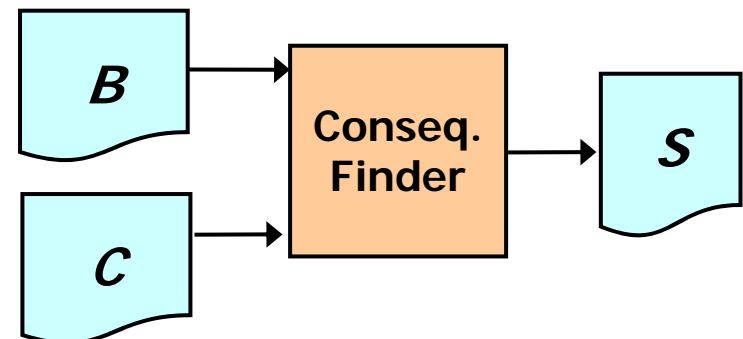
Consequence Finding

Input:

- B : first-order (clausal) theory
- C : “new” clausal theory
- P : language restriction (“*production field*”)

Output:

- S : the (minimal) “new” consequences satisfying that
 - $B \wedge C \models S$
 - $B \not\models S$
 - S belongs to P .
- **SOL-resolution (Inoue, IJCAI-91)**
- **SOLAR (Nabeshima, Iwanuma & Inoue, TABLEAUX’03)**
- For Theorem Proving, C is the negation of the target theorem and S is the empty clause (generalization of *proof-finding*).
- For Inverse Entailment, $C = \neg E$ and $S = \neg H$: $B \wedge \neg E \models \neg H$.



Inverse Entailment for Abduction

SOLAR Example: graph completion problem – *pathway finding*

Find an arc which enables a path from A to D.

Background theory:

$\text{path}(X,Y) \leftarrow \text{node}(X), \text{node}(Y), \text{arc}(X,Y).$

$\text{path}(X,Z) \leftarrow \text{node}(X), \text{node}(Y), \text{node}(Z), \text{arc}(X,Y), \text{path}(Y,Z).$

$\text{node}(a). \text{node}(b). \text{node}(c). \text{node}(d). \text{arc}(a,b). \text{arc}(c,d).$

Negated observation:

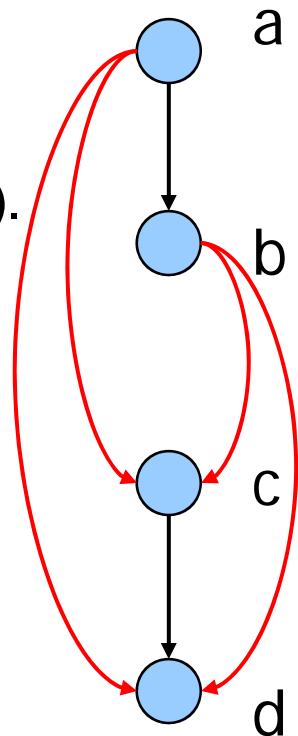
$\neg \text{path}(a,d).$

Production field:

literal form = $[\neg \text{arc}(_,_)]$ & clause length = 1.

Output of SOLAR:

1. $\neg \text{arc}(a, d).$
2. $\neg \text{arc}(a, c).$
3. $\neg \text{arc}(b, c).$
4. $\neg \text{arc}(b, d).$



Metabolic Pathway (Ray & Inoue, DS'07)

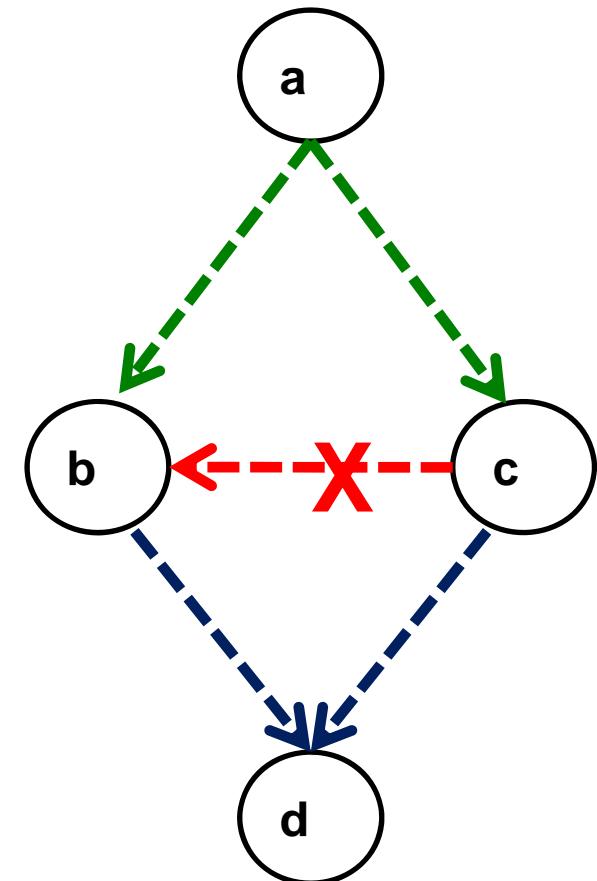
$$T = \begin{cases} pathway(X, Z) \leftarrow reaction(X, Y) \wedge pathway(Y, Z) \\ pathway(X, Z) \leftarrow reaction(X, Z) \\ reaction(a, b) \vee reaction(a, c) \\ reaction(b, d) \vee reaction(c, d) \\ \neg reaction(c, b) \end{cases}$$

$$E = \{pathway(U, d)\}$$

% (from which U) is there a path to d?

$$A = \{reaction(V, W)\}$$

% assuming reactions from some V to W



Metabolic Pathway: A Solution

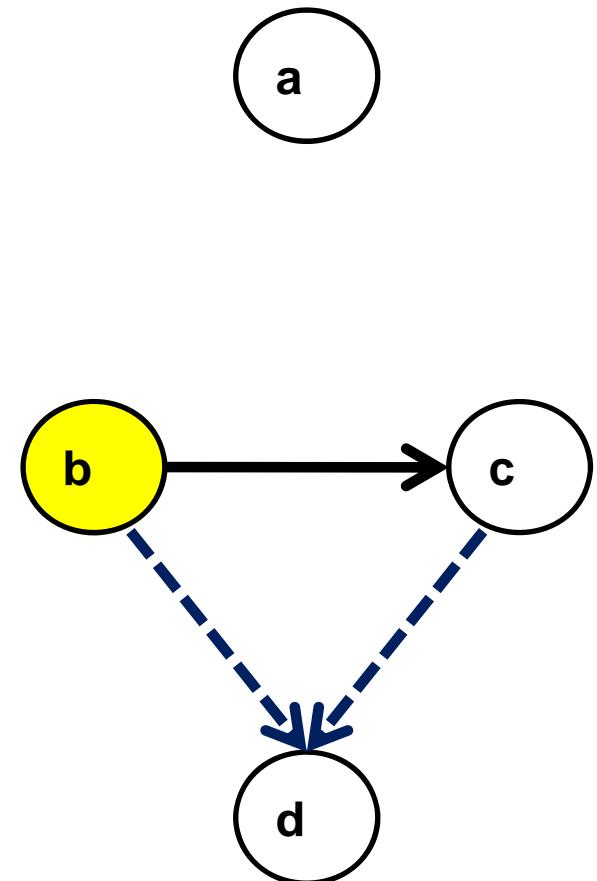
$$B = \begin{cases} pathway(X, Z) \leftarrow reaction(X, Y) \wedge pathway(Y, Z) \\ pathway(X, Z) \leftarrow reaction(X, Z) \\ reaction(a, b) \vee reaction(a, c) \\ reaction(b, d) \vee reaction(c, d) \\ \neg reaction(c, b) \end{cases}$$

$$E\sigma = \{pathway(b, d)\}$$

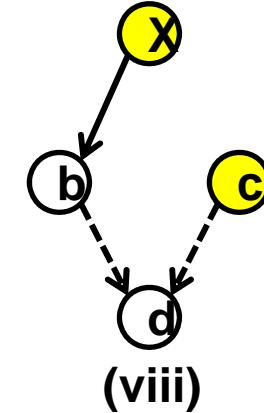
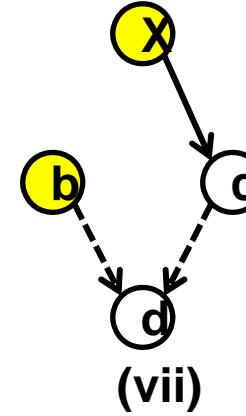
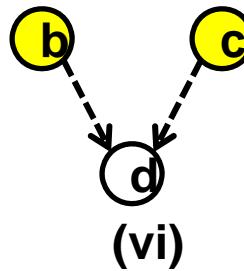
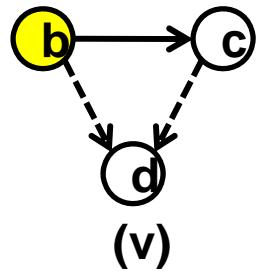
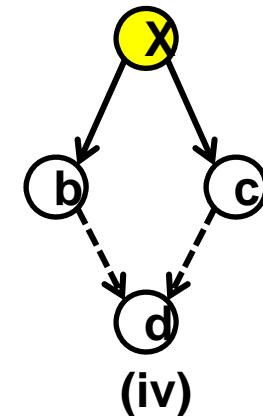
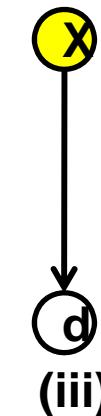
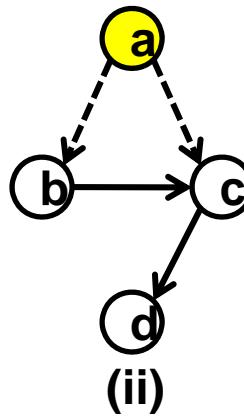
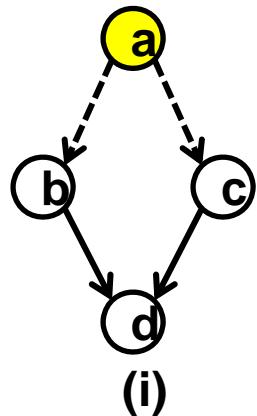
% there is a path from b to d

$$H = \{reaction(b, c)\}$$

% assuming a reaction from b to c



Metabolic Pathway: MORE Solutions



Problem: want to express non-ground answers like (iii), and disjunctive answers such as (vi).

Abductive Inference (Ray & Inoue, DS'07)

given

B

theory (background knowledge)

E

goal (set of given observations)

A

abducibles (set of possible assumptions)

find

$H \subseteq A$

explanation (set of assumptions)

Θ

answer (SET of variable bindings)

where

$$B \models \forall \left(\bigwedge_{L \in H} L \rightarrow \bigvee_{\sigma \in \Theta} E\sigma \right)$$

% the conjunction of assumptions H implies the disjunction of answers Θ .

Metabolic Pathway: Another Solution

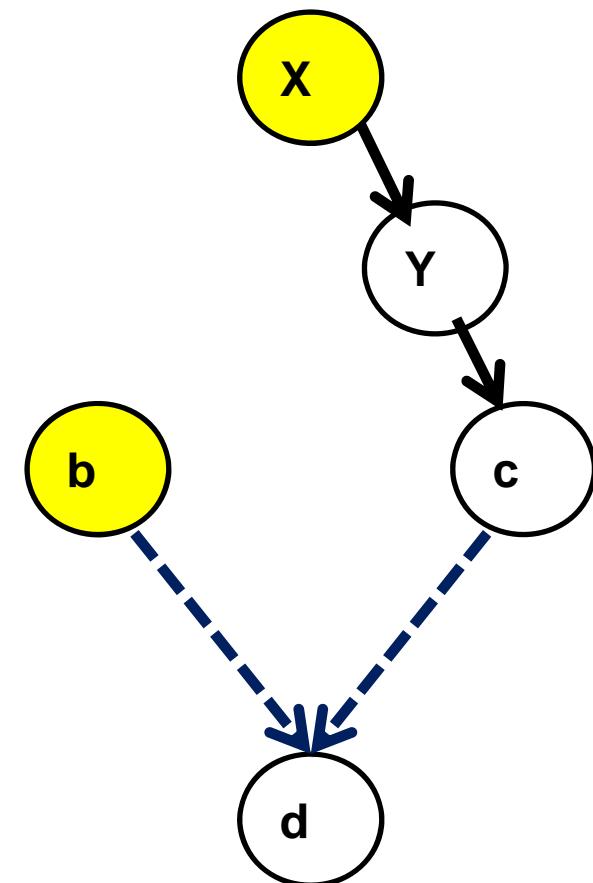
$$B = \begin{cases} pathway(X, Z) \leftarrow reaction(X, Y) \wedge pathway(Y, Z) \\ pathway(X, Z) \leftarrow reaction(X, Z) \\ reaction(a, b) \vee reaction(a, c) \\ reaction(b, d) \vee reaction(c, d) \\ \neg reaction(c, b) \end{cases}$$

$$\Theta = \{\{U / b\}, \{U / X\}\}$$

% there is a path from b or X to d

$$H = \{reaction(X, Y), reaction(Y, c)\}$$

% assuming reactions from X to Y and from Y to c

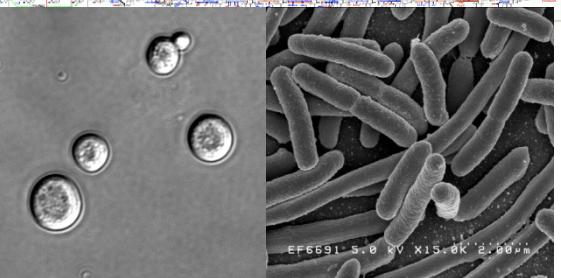
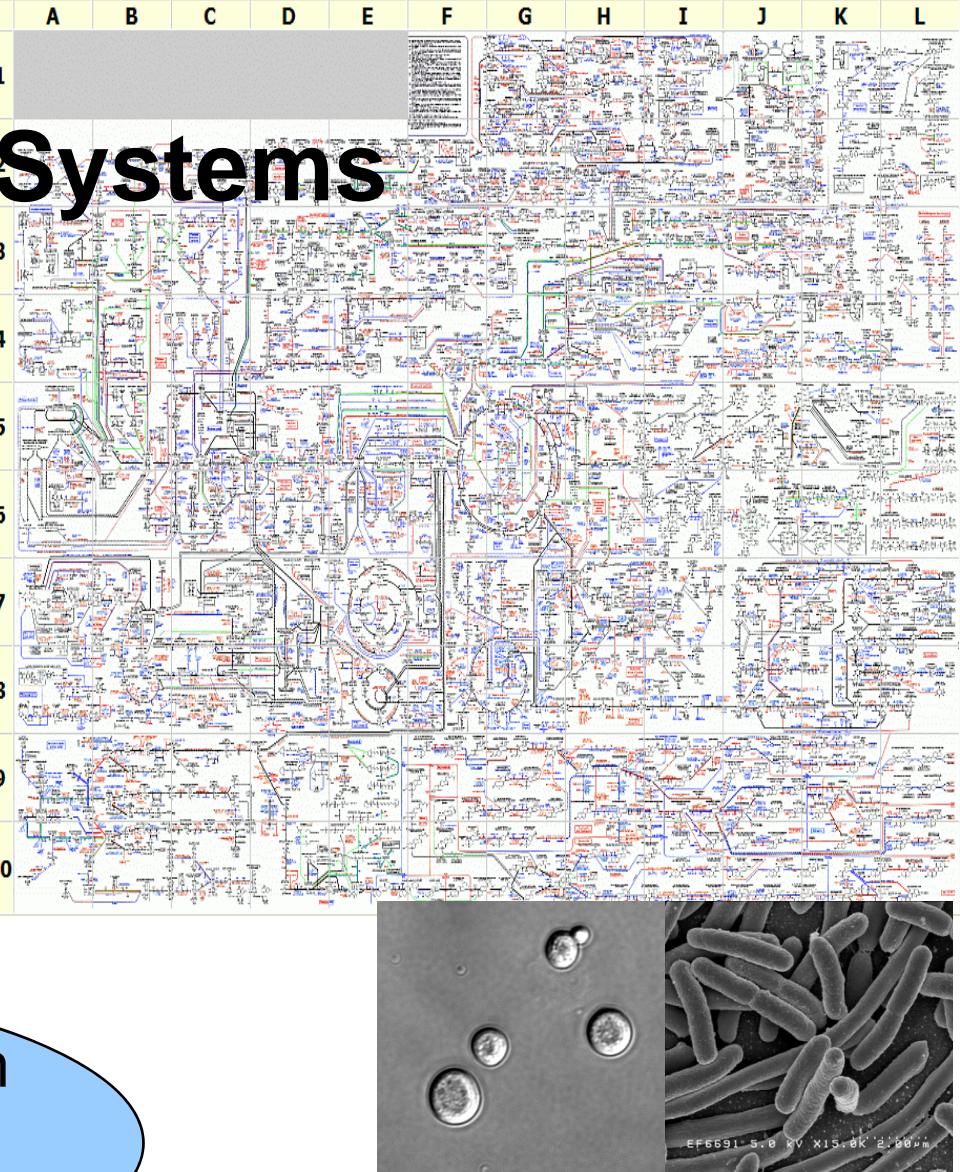


Modeling Biological Systems

- Explain and predict metabolic pathways.

- Generic Model:
 - *Saccharomyces Cerevisiae*
 - *E-coli*

- Biological Phenomenon can be explained by *Inductive Logic Programming (ILP)*.



Inductive Learning Approaches

- **Goals**

- Finding inhibitions in a metabolic pathway.
- Discovering causal rules which augment an incomplete background theory.
- Predicting changes of concentration in intracellular fluxes.

- **Previous Work**

- Using an **abductive logic programming** technique on the problem of inhibitions of metabolic pathways at steady states (Tamaddoni-Nezdah et al., 2006)

- **New Approach** (Yamamoto, Inoue & Doncescu, 2007)
 - Integration of **abduction** and **induction**.
 - Not only **steady states** but also **dynamic models**.

Inverse Entailment for Induction

- Horn clauses for background theories
 - **Progol** (Muggleton, 1995)
- Full clausal theories (non-Horn clauses) for background, example, and hypothesis theories.
 - **CF-induction** (Inoue, ILP-2001; Yamamoto, Ray & Inoue, LLLL-2007)
 - **fc-HAIL** (Ray & Inoue, ILP-2007)
- Note: CF-induction is the only existing ILP system that is **complete** for full clausal theories.

Example: Metabolic Pathway (Pyruvate)

■ **B:**

reaction(pyruvate, acetylcoa).

reaction(pyruvate, acetaldehyde).

reaction(glucose, glucosep).

reaction(glucosep, pyruvate).

reaction(acetaldehyde, acetate).

reaction(acetate, acetylcoa).

reaction(acetaldehyde, ethanol).

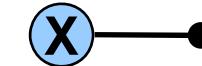
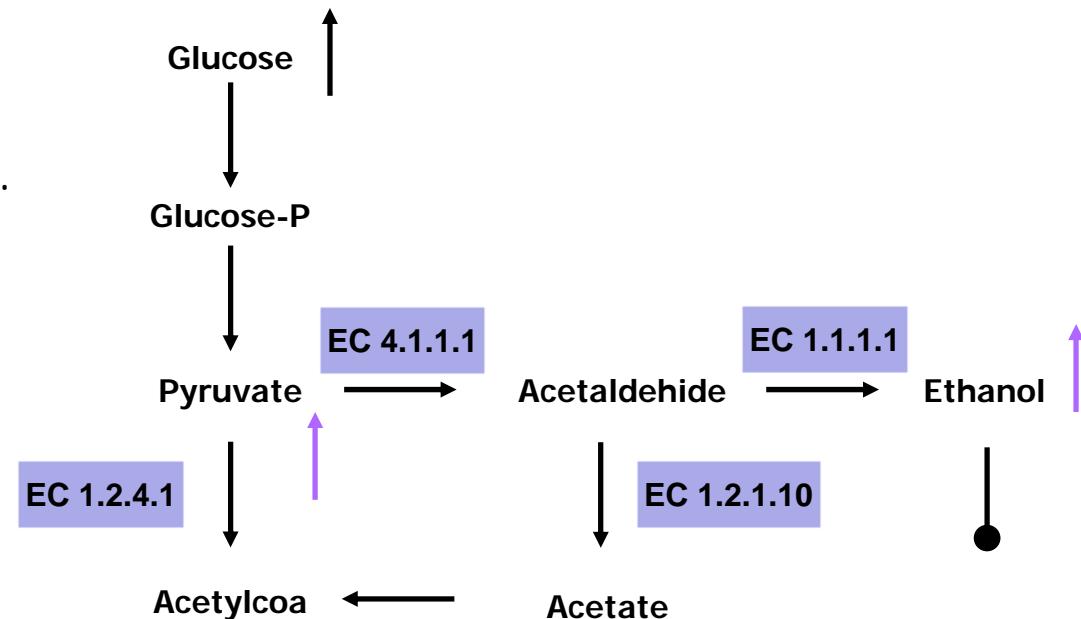
concentration(glucose, up).

terminal(ethanol).

blocked(X) \leftarrow reaction(X,Y), inhibited(X,Y).

blocked(X) \leftarrow terminal(X).

concentration(X,up) \leftarrow reaction(Y,X), \neg inhibited(Y,X), blocked(X).

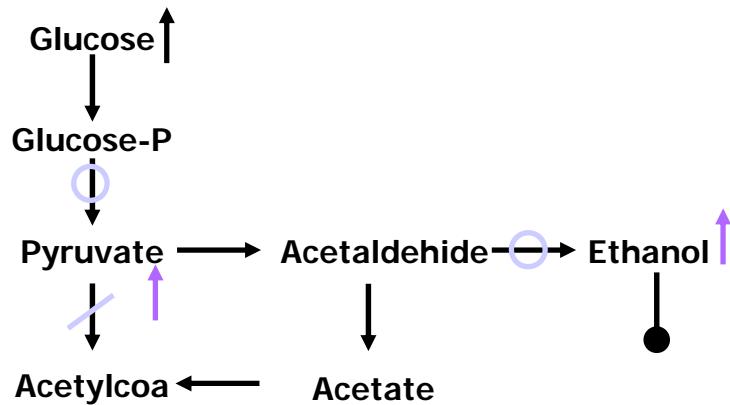


■ **E:** concentration(ethanol,up). concentration(pyruvate, up).

Example: Outputs of CF-induction

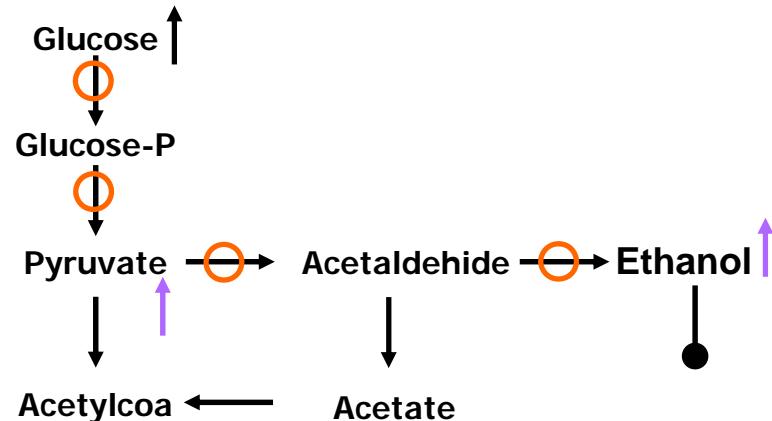
■ H_1 :

- ¬Inhibited(glucosep, pyruvate).
- ¬inhibited(acetaldehyde, ethanol).
- inhibited(pyruvate, acetylcoa).



■ H_2 :

- ¬inhibited(glucose, glucosep)
- ¬Inhibited(glucosep, pyruvate).
- ¬inhibited(acetaldehyde, ethanol).
- ¬inhibited(pyruvate, acetaldehyde).
- concentration(Y, up) ←
- ¬inhibited(X, Y), concentration(X, up).



Conclusion

- Algorithm = Logic + Control.
- Logic programming is still alive.
- Is this GOFAI?
- Dealing with **incomplete knowledge** and **elaboration tolerance** — default reasoning, planning, prediction, explanation, hypothesis formation, model generation.
- Inference-based hypothesis-finding in biochemical networks.
- Logic is a glue to combine mathematical/abstract models with the real world.
- Get back to logic (programming)!